CrashOS: Hypervisor testing tool

ISSRE 2017

Anaïs GANTET - Airbus Digital Security

October 2017



Outline

Why CrashOS?

CrashOS presentation

Vulnerability research and results



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Why CrashOS?

CrashOS presentation

3 Vulnerability research and results



Project context

Goal

Test the robustness of virtualization solutions

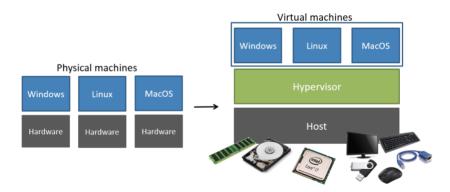
Scope

- Targeted architecture: Intel x86
- Initial targeted software: Ramooflax (https://github.com/airbus-seclab/ramooflax)
- Other targeted software: VMware, Xen, etc.



Virtualization in a nutshell

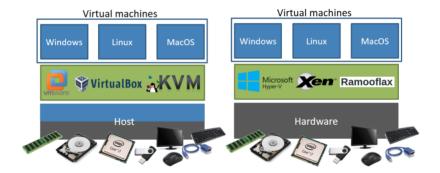
Aim: allow several Operating Systems (OS) to share hardware resources





Virtualization in a nutshell

Aim: allow several Operating Systems (OS) to share hardware resources





Hypervisor role

To provide an equivalent environment to a physical machine

- Processor virtualization
 - Let directly execute non-sensitive instructions
 - Trap and virtualize sensitive instructions
- Memory access virtualization
- Device virtualization

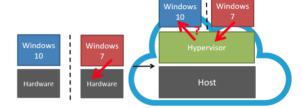
Hypervisors are

- Complex programs
- Operating at low level



Security issues

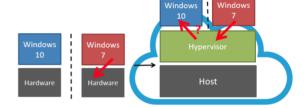
- Without virtualization: physical isolation
- With virtualization: software isolation (implemented by hypervisors)





Security issues

- Without virtualization: physical isolation
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Problem

Are hypervisors reliable?



Our approach

Behaviour analysis

- Build a testing VM
- Imagine tests involving the hypervisor
 - Be aware of known hypervisor weak points (CVE)
 - Instruction disassembly
 - Sensitive instruction emulation
 - Unusual CPU configuration
 - Device virtualization

 - etc.
- Launch them and observe hypervisor behaviour



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Our way to build a malicious VM

Potential candidates?

- Windows, Linux: only partial control of hardware communication
- Simple Operating System, OSv, etc.: not appropriate for vulnerability research
- VESPA: Only part targeting device virtualization is published





Our way to build a malicious VM

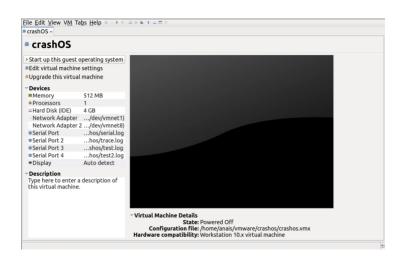
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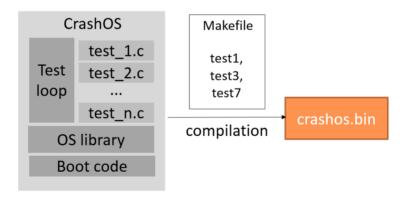
Our choice: develop our own OS dedicated to test hypervisors







How does CrashOS work?





CrashOS test format

- Init function
 - Save the current machine state
 - Define the desired context
- Test function
 - Execute instruction
 - Print adapted logs
- Restore function
 - Reinit the saved machine state

```
test_t test_x = {
    .name = "test name",
    .desc = "test description",
    .init = init_test_x,
    .test = test_test_x,
    .fini = restore_test_x
};
DECLARE_TEST(test_x)
```



CrashOS: Core and libraries

System features

- Physical memory access
- Memory protection mechanisms
- Interrupt handling
- Device communication
- Paravirtualization support

Example:

```
#define enable_paging()
  asm volatile (
  "mov %%cr0, %%eax \n"
  "or $0x80000000, %%eax \n"
  "mov %%eax, %%cr0":::"eax" )
```



CrashOS: Core and libraries

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```

Other interesting features

- Printing logs on the screen
- Printing logs on the serial port



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- 1 Why CrashOS?
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Attack approach

Prerequisites

- Advanced understanding of Intel processor mechanisms
- Good knowledge of hypervisor role

Elaboration of relevant tests

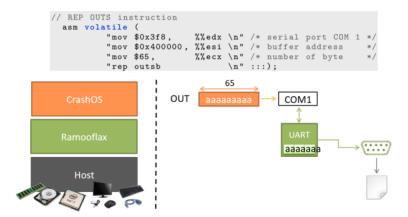
- Not only fuzz all instructions
- Always involve a specific hypervisor handling
 - Sensitive instructions
 - Memory management
 - Device emulation
 - etc.





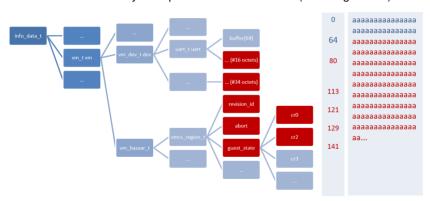
```
REP OUTS instruction
asm volatile (
       "mov $0x3f8, %%edx \n" /* serial port COM 1
       "mov $0x400000, %%esi \n" /* buffer address
       "mov $65.
                       %%ecx \n" /* number of byte
                              \n" :::);
       "rep outsb
                   OUT
                                        COM1
  Host
```







Buffer overflow and memory corruption of Ramooflax data (including VMCS)









Buffer overflow in Ramooflax causes a VMware crash



A second example: Privilege changing (FAR JMP)

Our test case

Launch FAR JMP from high level to low level (not allowed - segmentation fault must occur)

What is FAR JMP instruction?

- Changes the current execution context to a new one
 - Can modify the current privilege level

Why to test this instruction?

- FAR JMP success: only if parameters meet about 20 conditions
- Sensitive instruction, intercepted and handled by hypervisors
- Test based on CVE-2014-8595 Xen HVM



A second example: Privilege changing (FAR JMP)

Xen (HVM):

```
root@debian:~# xl list
                                                                   Time(s)
Name
                                             Mem VCPHs
                                                           State
Domain-0
                                                                    220.9
crashos
                                                                      5.7
root@debian:~# xl dmesg_
     domain crash called from vmx.c:2274
      Domain 28 (vcpu#0) crashed on cpu#0:
      ----[ Xen-4.2.4-pre x86 32n
                                      debug=n Not tainted ]----
(XEN)
(XEN)
              0030:[<00201b95>] EIP 2
(XEN)
              00000006
                          CONTEXT: hvm guest
(XEN)
      eax: 000h8000
                       ebx: 0002dae0
                                        ecx: 000b8000
                                                         edx: 000b8000
      esi: 00054769
                       edi:
                            0005476a
                                        ebp: 0020721c
      cr0: 00000011
                       cr4: 00000000
                                        cr3: 00800000
                                                         cr2: 0000
     ds: 0023
                     0000
                             fs: 0000
                                             0000
                                                     ss: 0023
```



A second example: Privilege changing (FAR JMP)

VMware:



CrashOS: Some test cases

	VMware	Ramooflax	Xen(HVM)
Access physical memory out of RAM size		3	9
Use memory with bad rights			
Mix 16-bits and 32-bits code		\sim	
Read or write on all I/O ports	\sim	3	
Badly configure the paging tables	\sim		
Badly emulate context switching		3	\sim

9	Correctly handled
	Not implemented
	VM crash (DoS)
3	More critical impacts



Results and discussions

- Vmware:
 - Monitor panic VERIFY
 - Monitor panic NOT IMPLEMENTED
 - Monitor panic NOT REACHED
 - Monitor panic EPT Misconfiguration

Contact: VMware Security Response Center

- Xen
 - Bad emulation of far imp and far call instructions reported
 - Ongoing to be fixed (https: //lists.xenproject.org/archives/html/xen-devel/2017-09/msg03701.html)
- Ramooflax
 - All bugs fixed



Conclusion and future outlook

CrashOS today

- Opensource
- Configurable minimal OS
- Simplified framework to write attacks on hypervisors
- First results on VMware, Xen and Ramooflax

Future outlook

- Elaborate new tests
- Implement new CrashOS features (64 bit support, *nested virtualization*, etc.)
- Test other hypervisors (Xen PV, KVM, Virtualbox, etc.)



Thanks for your attention

Questions?

https://github.com/airbus-seclab/crashos

French paper: https://www.sstic.org/media/SSTIC2017/SSTIC-actes/crashos/ SSTIC2017-Article-crashos-gantet.pdf

anais.gantet@airbus.com